

Limits of discrete structures and their applications

Dan Král'

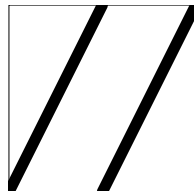


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Bedlewo, Poland

WHY LIMITS?

- asymptotic properties of large (discrete) objects
we implicitly use limits in our considerations anyway
- How does the seq. $1, 3, \dots, 2n - 1, 2, 4, \dots, 2n$ look like?
How does the adjacency matrix of $K_{n,n}$ look like?
How does the adjacency matrix of $K_{n,n+1}$ look like?
- convergence of a sequence of discrete objects vs.
formal analytic representation of its limit



$$\begin{array}{|c|c|} \hline 0 & 1 \\ \hline 1 & 0 \\ \hline \end{array}$$

OVERVIEW OF THE TALK

- Limits of dense structures
Survey of main results in the area
Limits of permutations and dense graphs
- The flag algebra method
Applications in extremal combinatorics
- Limits of sparse structures
Various concepts, less understood

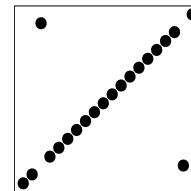
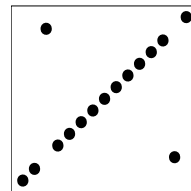
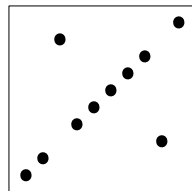
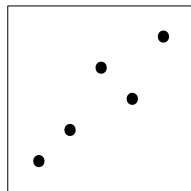
PERMUTATIONS

- permutation of order n : order on numbers $1, \dots, n$
subpermutation: $4\underline{53}21\underline{6} \longrightarrow 213$

- density of a permutation π in a permutation Π :

$$d(\pi, \Pi) = \frac{\# \text{ subpermutations of } \Pi \text{ that are } \pi}{\# \text{ all subpermutations of order } \pi}$$

- $(\Pi_j)_{j \in \mathbb{N}}$ convergent if $\exists \lim_{j \rightarrow \infty} d(\pi, \Pi_j)$ for every π



REPRESENTATION OF A LIMIT

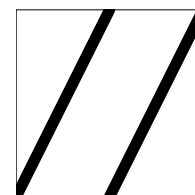
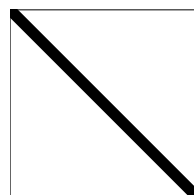
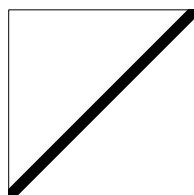
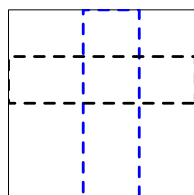
- probability measure μ on $[0, 1]^2$ with unit marginals

$$\mu([a, b] \times [0, 1]) = \mu([0, 1] \times [a, b]) = b - a$$

Hoppen, Kohayakawa, Moreira, Ráth and Sampaio
similar ideas in work of Presutti and Stromquist

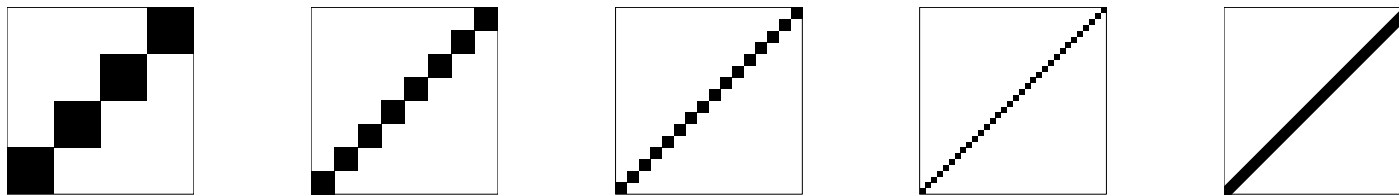
- μ -random permutation

choose n random points, x - and y -coordinates



EXISTENCE AND UNIQUENESS

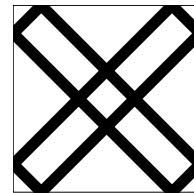
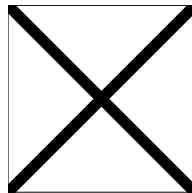
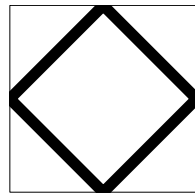
- **existence:** associate Π_j with a measure μ_j
 observe $|d(\pi, \Pi_j) - d(\pi, \mu_j)| \leq O(|\pi|^2 \cdot |\Pi_j|^{-1})$
 set $\mu(A) := \lim_{j \rightarrow \infty} \mu_j(A)$ for $A \subseteq [0, 1]^2$



- **uniqueness:** $\mu(A) = \lim_{n \rightarrow \infty} \mathbb{E} \frac{|\{m, [\frac{m}{2^n}, \frac{a_m}{2^n}] \in A\}|}{2^n}$
 where a_1, \dots, a_{2^n} is μ -random and $A \subseteq [0, 1]^2$
 $\forall \pi \ d(\pi, \mu) = d(\pi, \mu') \implies \mu = \mu'$

APPLICATION: QUASIRANDOMNESS

- property $P(k)$ of a sequence $(\Pi_j)_{j \in \mathbb{N}}$:
 $d(\pi, \Pi_j) \rightarrow 1/k!$ for every $\pi \in S_k$
- Question (Graham): Is there k_0 such $\forall k P(k_0) \Rightarrow P(k)$?
- Theorem (K., Pikhurko): yes, $k_0 = 4$; best possible
density of 123 is $1/8$ in the left and $1/4$ in the middle
a μ -random permutation based on the right measure

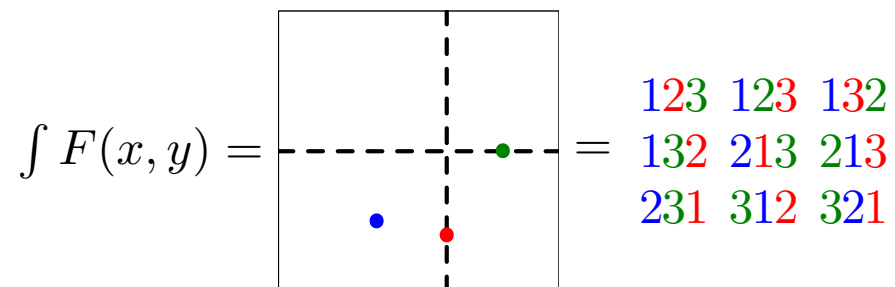
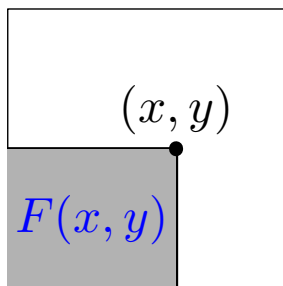


PROOF $P(5) \Rightarrow$ QUASIRANDOMNESS

- relate integrals and permutation densities

$$\int x \, dx dy = \frac{d(12, \mu) + d(21, \mu)}{2} = \frac{1}{2} = \int x \, dx$$

prob. $x_2 \leq x_1$ for two random points $[x_1, y_1]$ and $[x_2, y_2]$



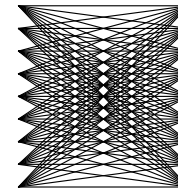
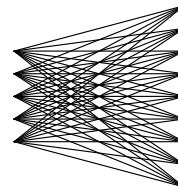
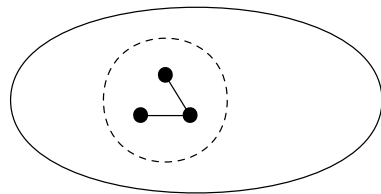
$$\int F(x, y) \, dx dy = \frac{d(123) + d(132) + d(213)}{3} + \frac{d(231) + d(312) + d(321)}{6}$$

$$\frac{1}{81} = \left(\int F(x, y) xy \, dx dy \right)^2 \leq \frac{1}{9} \int F(x, y)^2 dx dy = \frac{1}{81}$$

Questions?

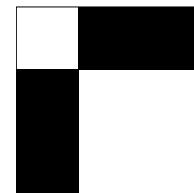
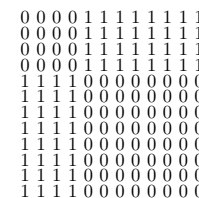
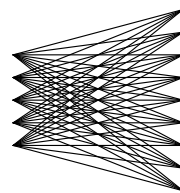
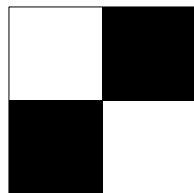
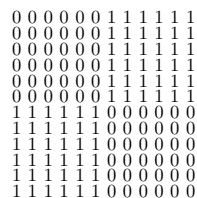
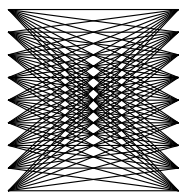
DENSE GRAPH CONVERGENCE

- Borgs, Chayes, Lovász, Sós, Szegedy and Vesztergombi
- convergence for **dense** graphs ($|E| = \Omega(|V|^2)$)
- $d(H, G) =$ probability $|H|$ -vertex subgraph of G is H
- a sequence $(G_n)_{n \in \mathbb{N}}$ of graphs is L-convergent if $d(H, G_n)$ converges for every H
- extendable to other discrete structures



LIMIT OBJECT: GRAPHON

- graphon $W : [0, 1]^2 \rightarrow [0, 1]$, s.t. $W(x, y) = W(y, x)$
- W -random graph of order n
random points $x_i \in [0, 1]$, edge probability $W(x_i, x_j)$
- $d(H, W) = \text{prob. } |H|\text{-vertex } W\text{-random graph is } H$
- W is a limit of $(G_n)_{n \in \mathbb{N}}$ if $d(H, W) = \lim_{n \rightarrow \infty} d(H, G_n)$



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- W is a limit of $(G_n)_{n \in \mathbb{N}}$ if $d(H, W) = \lim_{n \rightarrow \infty} d(H, G_n)$
- every L-convergent sequence of graphs has a limit
- W -random graphs converge to W with probability one

W -RANDOM GRAPHS

- the density of a graph H in a graphon W :

$$\frac{|H|!}{|\text{Aut}(H)|} \int_{[0,1]^{|H|}} \prod_{v_i v_j} W(x_i, x_j) \prod_{\frac{v_i v_j}{v_i v_j}} (1 - W(x_i, x_j)) dx_1 \cdots x_n$$

- a sequence $(G_n)_{n \in \mathbb{N}}$ of W -random graphs, $|G_n| = n$

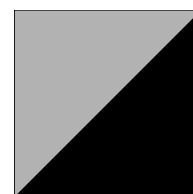
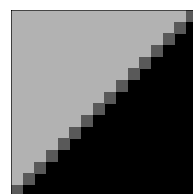
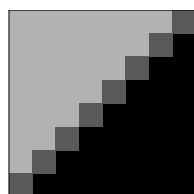
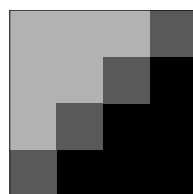
the expectation of $d(H, G_n)$ conditioned on x_1, \dots, x_i

Azuma's ineq.: $\mathbb{P}[|d(H, G_n) - d(H, W)| \geq \varepsilon] \leq e^{-O(\varepsilon^2 \cdot n)}$

Borel-Cantelli $\Rightarrow (G_n)_{n \in \mathbb{N}}$ converges with prob. one

CONSTRUCTION OF THE LIMIT

- sequence of mutually refining regularity partitions
removal lemma \Rightarrow subgraph counts
- interpret the partitions as functions $[0, 1]^2 \rightarrow [0, 1]$
the pointwise limit is the sought graphon
existence by martingale convergence

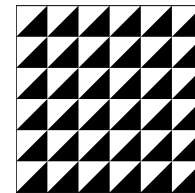
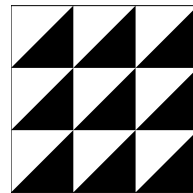
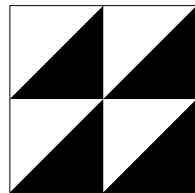


UNIQUENESS OF THE LIMIT

- $W^\varphi(x, y) := W(\varphi(x), \varphi(y))$ for $\varphi : [0, 1] \rightarrow [0, 1]$
- $d(H, W) = d(H, W^\varphi)$ if φ is measure preserving
- Theorem (Borgs, Chayes, Lovász)

If $d(H, W_1) = d(H, W_2)$ for all graphs H ,

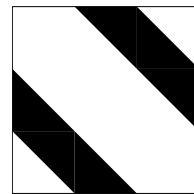
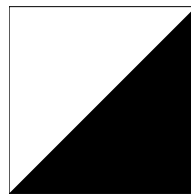
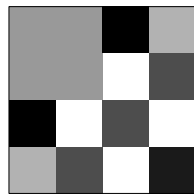
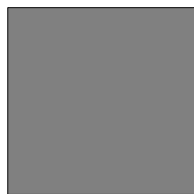
then there exist measure preserving maps φ_1 and φ_2 such that $W_1^{\varphi_1} = W_2^{\varphi_2}$ almost everywhere.



SPACE OF TYPICAL VERTICES

- vertices of a graphon W : its rows $f_x(y) := W(x, y)$
 $R(W) = \{f \in L_1[0, 1] \mid \exists x f = f_x\}$
 $R(W)$ is naturally equipped with measure μ
- **typical vertices** $T(W)$ = the support of μ in $L_1[0, 1]$
topology can be given by L_1 or by similarity distance

$$d_W(f, f') = \mathbb{E}_{g_x} \left| \int_y (f(y) - f'(y)) g_x(y) \right|$$



WEAK REGULARITY PARTITIONS

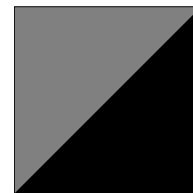
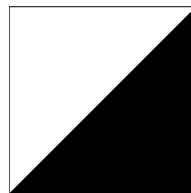
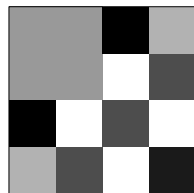
- Frieze and Kannan (1999)
an ε -regular partition of a graph G is V_1, \dots, V_k

$$\left| e(S, T) - \sum_{i,j} d_{ij} \cdot |S \cap V_i| \cdot |T \cap V_j| \right| \leq \varepsilon$$

- \exists ε -regular partitions with k parts \Leftrightarrow
 \exists cover of typical vertices with k ε -balls except for ε
- finite dimension \Rightarrow polynomial weak regularity partition

FINITELY FORCIBLE GRAPHONS

- a graphon W is **finitely forcible** if there exist H_1, \dots, H_k and d_1, \dots, d_k such that W is the only graphon with the expected density of H_i equal to d_i
- \Leftrightarrow the only graphon minimizing $\sum \alpha_j d(H'_j, W)$
- Lovász and Sós (2008),
Every step function is finitely forcible.



FINITELY FORCIBLE GRAPHONS

- Conjecture (Lovász and Szegedy):
Every extremal problem $\min \sum \alpha_j d(H_j, W)$
has a finitely forcible optimal solution.
- extremal graph theory problem \rightarrow
finitely forcible optimal solution \rightarrow
simple structure gives new bounds on old problems
(bounded dimension, few kinds of typical vertices)

FINITELY FORCIBLE GRAPHONS

- Conjectures (Lovász and Szegedy):
The space $T(W)$ of a finitely forcible W is compact.
The space $T(W)$ has finite dimension.
- Theorem (Glebov, K., Volec):
 $T(W)$ can fail to be locally compact
- Theorem (Glebov, Klimošová, K.):
 $T(W)$ can have a part homeomorphic to $[0, 1]^\infty$
- Theorem (Cooper, Kaiser, K., Noel):
 \exists finitely forcible W such that every ε -regular partition
has at least $2^{\varepsilon^{-2}} / \log \log \varepsilon^{-1}$ parts.

PARAMETER TESTING

- graph parameter $\mathcal{P} : \text{graphs} \rightarrow \mathbb{R}$
- large input data, not possible to process
providing an estimate based on a small sample
- \mathcal{P} is **testable** if there exists a randomized algorithm that estimates the parameter \mathcal{P} within the **additive error** ε based on a **sample of size** $f(\varepsilon)$ with **probability** $\geq 1 - \varepsilon$
- \mathcal{P} is testable $\Leftrightarrow \mathcal{P}$ is continuous on the graphon space

Questions?

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FLAG ALGEBRAS: HOMOMORPHISMS

- algebra \mathcal{A} on formal linear combinations of graphs
addition and multiplication by a scalar
- if W is a graphon, define $f(\sum \alpha_i H_i) := \sum \alpha_i d(H_i, W)$
the function f is a **homomorphism** from \mathcal{A} to \mathbb{R}
we define multiplication in \mathcal{A} on the next slide
- $\text{Ker}(f)$ always contains certain elements
e.g. $K_2 - \frac{1}{3}\overline{K_{1,2}} - \frac{2}{3}K_{1,2} - \frac{3}{3}K_3$
 $f(K_2) = \frac{1}{3}f(\overline{K_{1,2}}) + \frac{2}{3}f(K_{1,2}) + \frac{3}{3}f(K_3)$

FLAG ALGEBRAS: MULTIPLICATION

- picking two pairs of vertices \approx picking a quadruple assuming the considered graph G is huge
- $d(K_2, G) \times d(K_2, G) = \frac{1}{3}d(K_2 \cup K_2, G) + \frac{1}{3}d(P_4, G) + \dots$
- this leads to a definition of a **product of two graphs**

$$\begin{array}{c} \bullet \\ | \\ \bullet \end{array} \times \begin{array}{c} \bullet \\ | \\ \bullet \end{array} = \frac{1}{3} \begin{array}{c} \bullet \\ | \\ \bullet \end{array} + \frac{1}{3} \begin{array}{c} \bullet & & \bullet \\ | & & | \\ \bullet & \text{---} & \bullet \end{array} + \frac{2}{3} \begin{array}{c} \bullet & & \bullet \\ | & & | \\ \bullet & \text{---} & \bullet \\ | & & | \\ \bullet & & \bullet \end{array} + \frac{1}{3} \begin{array}{c} \bullet & & \bullet \\ | & & | \\ \bullet & \text{---} & \bullet \\ & \diagup & \\ & \bullet & \\ & \diagdown & \\ \bullet & & \bullet \end{array} + \frac{2}{3} \begin{array}{c} \bullet & & \bullet \\ | & & | \\ \bullet & \text{---} & \bullet \\ & \diagup & \\ & \bullet & \\ & \diagdown & \\ \bullet & & \bullet \\ & \diagup & \\ & \bullet & \\ & \diagdown & \\ \bullet & & \bullet \end{array} + 1 \begin{array}{c} \bullet & & \bullet \\ | & & | \\ \bullet & \text{---} & \bullet \\ & \diagup & \\ & \bullet & \\ & \diagdown & \\ \bullet & & \bullet \\ & \diagup & \\ & \bullet & \\ & \diagdown & \\ \bullet & & \bullet \end{array}$$

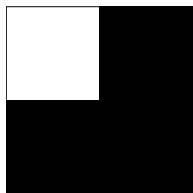
- If W is the unique graphon with $d(H_i, W) = \alpha_i$,

it is the unique graphon with $\sum_i (d(H_i, W) - \alpha_i)^2 = 0$.

$$\sum_i (H_i - \alpha_i K_1)^2 = \sum_j \beta_j H'_j$$

ROOTED HOMOMORPHISMS

- consider a graph G with a distinguish vertex (root)
a random sample always includes the root
- algebra \mathcal{A}^\bullet on combinations of rooted graphs
- rooted graph \rightarrow a homomorphism from \mathcal{A}^\bullet to \mathbb{R}
random choice of the root \rightarrow probability distribution
on homomorphisms f^\bullet from \mathcal{A}^\bullet to \mathbb{R}



$$f^\bullet(K_2^\bullet) = 1/2, f^\bullet(\overline{K_2^\bullet}) = 1/2, f^\bullet(K_3^\bullet) = 0, \dots$$

$$f^\bullet(K_2^\bullet) = 1, f^\bullet(\overline{K_2^\bullet}) = 0, f^\bullet(K_3^\bullet) = 3/4, \dots$$

AVERAGING OPERATOR

- goal: $f(\llbracket x \rrbracket_{\bullet}) = \mathbb{E}_{\bullet} f^{\bullet}(x)$ for $x \in \mathcal{A}^{\bullet}$

$$\llbracket \begin{array}{c} \circ \\ \diagdown \quad \diagup \\ \bullet \end{array} \rrbracket_{\bullet} = \frac{1}{3} \begin{array}{c} \circ \\ \diagdown \quad \diagup \\ \circ \end{array}$$

$$\llbracket \begin{array}{c} \circ \text{---} \circ \\ \diagdown \quad \diagup \\ \bullet \end{array} \rrbracket_{\bullet} = \frac{2}{3} \begin{array}{c} \circ \text{---} \circ \\ \diagdown \quad \diagup \\ \circ \end{array}$$

- expressing degree constraints:

$$f(\llbracket (K_2^{\bullet} - 1/3)^2 (K_2^{\bullet} - 2/3)^2 \rrbracket_{\bullet}) = 0$$

- if A is PSD $n \times n$ -matrix and $x \in \mathcal{A}^n$, $f(x^T A x) \geq 0$

$$\text{example: } f((\alpha K_2 - \beta \overline{K_2})^2) = f(\sum \alpha_i H_i) \geq 0$$

analogously for vectors x with entries from \mathcal{A}^{\bullet}

- search for such inequalities can be computer assisted

FLAG ALGEBRAS: EXAMPLE

$$\begin{array}{c}
 \begin{array}{c} \circ \\ \diagup \\ \bullet \end{array} \times \begin{array}{c} \circ \\ \diagup \\ \bullet \end{array} = \begin{array}{c} \circ \text{---} \circ \\ \diagup \quad \diagdown \\ \bullet \end{array} + \begin{array}{c} \circ \quad \circ \\ \diagdown \quad \diagup \\ \bullet \end{array} \\
 \begin{array}{c} \circ \\ \bullet \end{array} \times \begin{array}{c} \circ \\ \bullet \end{array} = \begin{array}{c} \circ \text{---} \circ \\ \bullet \end{array} + \begin{array}{c} \circ \quad \circ \\ \bullet \end{array}
 \end{array}$$

$$\begin{array}{c}
 \left[\left(\begin{array}{c} \circ \\ \diagup \\ \bullet \end{array} - \begin{array}{c} \circ \\ \bullet \end{array} \right)^2 \right]_{\bullet} = \frac{3}{3} \begin{array}{c} \circ \text{---} \circ \\ \diagup \quad \diagdown \\ \bullet \end{array} - \frac{1}{3} \begin{array}{c} \circ \quad \circ \\ \diagdown \quad \diagup \\ \bullet \end{array} - \frac{1}{3} \begin{array}{c} \circ \text{---} \circ \\ \circ \end{array} + \frac{3}{3} \begin{array}{c} \circ \quad \circ \\ \circ \end{array} \geq 0 \\
 \frac{1}{3} \begin{array}{c} \circ \text{---} \circ \\ \diagup \quad \diagdown \\ \bullet \end{array} + \frac{1}{3} \begin{array}{c} \circ \quad \circ \\ \diagdown \quad \diagup \\ \bullet \end{array} + \frac{1}{3} \begin{array}{c} \circ \text{---} \circ \\ \circ \end{array} + \frac{1}{3} \begin{array}{c} \circ \quad \circ \\ \circ \end{array} = \frac{1}{3} \\
 \begin{array}{c} \circ \text{---} \circ \\ \diagup \quad \diagdown \\ \bullet \end{array} + \begin{array}{c} \circ \quad \circ \\ \circ \end{array} \geq \frac{1}{4}
 \end{array}$$

QUASIRANDOMNESS OF GRAPHS

- Theorem (Thomason):

$$d_{\text{hom}}(H, G) = d(H, G_{n,p}) \text{ for } H = K_2, C_4 \Rightarrow G \approx G_{n,p}$$

- $d(H, W) = d(H, W_p)$ for all 4-vertex $H \Rightarrow W = W_p$

$$\forall x : \int W(x, y) \, dy = p \iff \llbracket (K_2^\bullet - p)^2 \rrbracket_\bullet = 0$$

$$\begin{aligned} \forall x, x' : \int W(x, y)W(x', y) \, dy &= p^2 \\ \iff \llbracket (K_3^{\bullet\bullet} + K_{1,2}^{\bullet\bullet} - p^2)^2 \rrbracket_{\bullet\bullet} &= 0 \\ \implies \forall x : \int W(x, y)^2 \, dy &= p^2 \end{aligned}$$

$$\begin{aligned} \forall x : \int W(x, y) \, dy = p \wedge \int W(x, y)^2 \, dy &= p^2 \\ \implies \forall x, y : W(x, y) &= p \end{aligned}$$

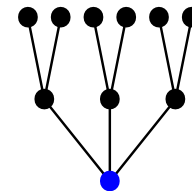
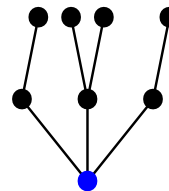
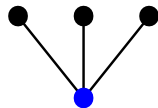
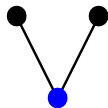
Questions?

SPARSE GRAPH CONVERGENCE

- convergence of graphs with bounded degree
trivially converging to the zero graphon
- need of a different notion of convergence
several notions, each having some cons
- absence of understood analytic representation
Aldous and Lyons Conjecture, relation to group theory
Does every graphing has a sequence converging to it?

LEFT CONVERGENCE

- introduced by Benjamini and Schramm in 2001
- bounded number of types of d -neighborhoods
convergence of statistic of d -neighborhoods
- cons: connected vs. disconnected
bipartite vs. non-bipartite graphs



LOCAL-GLOBAL CONVERGENCE

- introduced by Hatami, Lovász and Szegedy in 2012
- types of d -neighborhoods k -vertex-colored graphs
convergence of statistic of such d -neighborhoods
attainable by a k -vertex-coloring of graphs
- $(G_i)_{i \in \mathbb{N}} \rightarrow (A_i)_{i \in \mathbb{N}}$ where $A_i \subseteq \mathbb{R}^K$ and K is # of types
 $\forall \varepsilon > 0 \exists n \forall i, j > n, x \in A_i \exists y \in A_j \|x - y\| \leq \varepsilon$
- almost bipartite vs. non-bipartite graphs
local-global convergence \Rightarrow left convergence

GRAPHINGS

- graphing G is a graph with $V(G) = [0, 1]$
bounded maximum degree, Borel measurable edge-set
mass preservation: $\int_A \deg_B(x) dx = \int_B \deg_A(y) dy$
- Theorem (Elek, 2007)
Every BS-convergent sequence has a graphing.
Theorem (Hatami, Lovász, Szegedy, 2012)
Every LG-convergent sequence has a graphing.
- Conjecture (Aldous, Lyons)
Every graphing is a BS-limit of a graph sequence.

RIGHT CONVERGENCE

- developed by Borgs, Chayes, Kahn and Lovász in 2013

$$\lim_{i \rightarrow \infty} \frac{\log \text{hom}(G_i, H)}{|G_i|}$$

$$\text{hom}(G, H) = \sum_{f: G \rightarrow H} \prod_{v \in V} w(f(v)) \prod_{vv' \in E} w(f(v)f(v'))$$

- cons: connected vs. disconnected

$$\frac{\log \text{hom}(G, H)}{|G|} = \frac{\log \text{hom}(G \cup G, H)}{|G \cup G|}$$

RIGHT \Rightarrow LEFT CONVERGENCE

- complex proof by Borgs, Chayes, Kahn and Lovász
- use LLL to count $\text{hom}(G, H_k)$ for $H_k = K_k \setminus K_2$

if $p \leq x_i \prod_{j \sim i} (1 - x_j)$, then $\geq \prod_i (1 - x_i)$

$$p = \frac{2}{k^2} \leq \left(\frac{2}{k^2} + \frac{c}{k^4} \right) \left(1 - \frac{2}{k^2} - \frac{c}{k^4} \right)^{2\Delta}$$

$$\text{hom}(G, H_k) \geq k^{|G|} \cdot \left(1 - \frac{2}{k^2} - \frac{c}{k^4} \right)^{|E(G)|}$$

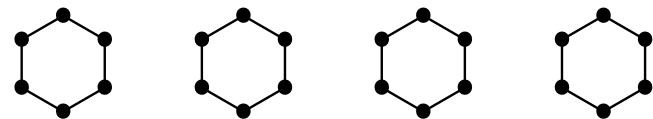
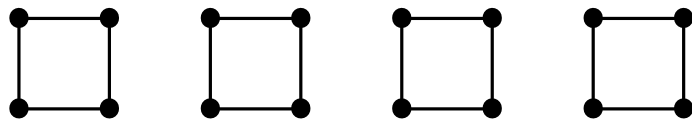
- upper bound given by Suen's inequality

$$\text{hom}(G, H_k) \leq k^{|G|} \cdot e^{\frac{2|E(G)|\Delta}{k^3}} \left(1 - \frac{2}{k^2} \right)^{|E(G)|}$$

- $\frac{\log \text{hom}(G, H_k)}{|G|} = k - \frac{|E(G)|}{|G| \cdot k^2} + O\left(\frac{\Delta^2}{k^3}\right)$

PARTITION CONVERGENCE

- introduced by Bollobás and Riordan in 2011
- statistic of a k -partition: $(a_1, \dots, a_k, d_{11}, d_{12}, \dots, d_{kk})$
convergence of attainable statistic of k -partitions
- $n \times C_4$ vs. $n \times C_6$: same attainable statistic
if $k = 2$: $a_1 + a_2 = 1$, $a_1 = d_{11} + d_{12}/2$, $a_2 = d_{22} + d_{12}/2$



LARGE DEVIATION CONVERGENCE

- introduced by Borgs, Chayes and Gamarnik in 2013
- counting k -partitions with a statistic $x \in \mathbb{R}^{k + \binom{k+1}{2}}$

$$r(x) = \lim_{\varepsilon \rightarrow 0} \lim_{i \rightarrow \infty} - \frac{\log \frac{\# \text{ } k\text{-partitions } \varepsilon\text{-close to } x}{k^{|G_i|}}}{|G_i|}$$

possible values: $[0, \log k] \cup \{\infty\}$

number of k -partitions with statistic $x \approx k^{|G|} \cdot e^{-r(x)|G|}$

- \Rightarrow partition convergence (the limit is finite)
- \Rightarrow right convergence (weighted partitions)

LARGE DEVIATION \Rightarrow RIGHT CONVERG.

- $\text{hom}(G, H) = \sum_{f:G \rightarrow H} \prod_{v \in V} w(f(v)) \prod_{vv' \in E} w(f(v)f(v'))$

- if $k = |H|$, determined by the statistic $x \in \mathbb{R}^{k + \binom{k+1}{2}}$

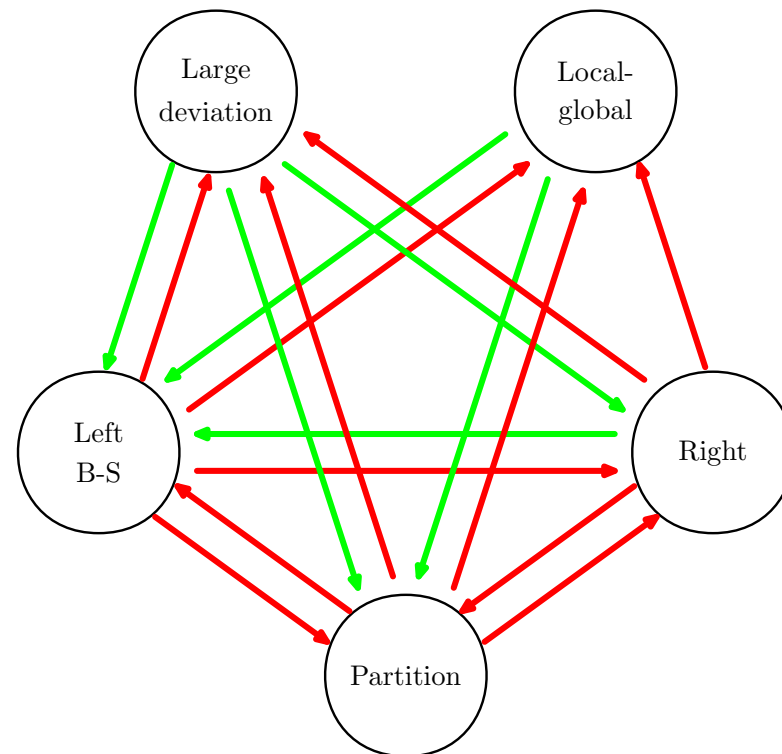
$$\approx \sum_x k^{|G|} e^{-r(x)|G|} \prod_{v_i} w(v_i)^{x_i|G|} \prod_{v_i v_j} w(v_i v_j)^{x_{ij}|G|}$$

- What is $\lim_{i \rightarrow \infty} \frac{\log \text{hom}(G_i, H)}{|G_i|}$?

$$\sup_x \log k - r(x) + \sum_{v_i} x_i \log w(v_i) + \sum_{v_i v_j} x_i x_j \log w(v_i v_j)$$

convergence of $r(x) \Rightarrow$ right convergence

MUTUAL RELATIONS



Questions?

EXERCISES

- Compute the density of 123 in the permutations below.
- Describe the limit of $K_{n/3} \cup K_{2n/3}$ and show that it is finitely forcible.
- Show that every BS-convergent sequence of 2-regular graphs has a graphing.

